

# Announcements

Tuesday, February 27

**Written Assignment:** *Prioritize quality over quantity.*

**Setting:** Representing a star-up, you are writing a report to a big company to offer your services. **In 2-3 pages, cover the following issues:**

- ▶ Give a 1/2 page summary of the proposal targeted to a layman
- ▶ What is the problem of the company (be creative)
- ▶ What is your proposed solution
- ▶ *Why does your solution works* (technical background)
- ▶ How will you implement the solution

Suggested companies/projects:

- ▶ Pixar's 3D simulations, Urban Planning, Tracing the trayjectory of Falcon 9

## Grading Scheme:

1. Correctness: 4pts
2. Clarity: 4pts
3. Creativity: 3pts

**Deadline:** Emailed in PDF by April 20th

## Group Sizes: 3-4 persons

- ▶ Suggested sources: Sections
- ▶ 2.6,2.7,3.3,4.8,4.9,5.6,5.7,5.8, 6.4,6.6,6.8,7.5,10.1,10.3,10.4
- ▶ *Not exclusively* Chapter 1

# Application

Stochastic Matrices and PageRank

# Stochastic Matrices

These arise very commonly in modeling of probabilistic phenomena (Markov chains), where they are also called **transition matrices**.

Some examples:

- ▶ Matrices from the population dynamics
- ▶ Matrices from the equilibrium-prices economies

## Definition

A square matrix  $A$  is **stochastic** if all of its entries are nonnegative, and the sum of the entries of each column is 1.

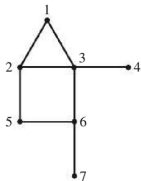
We say  $A$  is **regular** if, for some  $k$ , all entries of  $A^k$  are positive.

## Definition

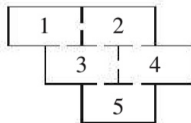
A **steady-state vector**  $v$  of  $A$  is a non-zero vector with *entries summing to 1* and such that  $Av = v$ .

# Random walks on graphs (a.k.a Mouse on a maze)

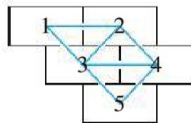
A mouse moves freely between rooms/states = selects any with equal probability.



$$P = \begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ 0 & 1/3 & 1/4 & 0 & 0 & 0 & 0 \\ 1/2 & 0 & 1/4 & 0 & 1/2 & 0 & 0 \\ 1/2 & 1/3 & 0 & 1 & 0 & 1/3 & 0 \\ 0 & 0 & 1/4 & 0 & 0 & 0 & 0 \\ 0 & 1/3 & 0 & 0 & 0 & 1/3 & 0 \\ 0 & 0 & 1/4 & 0 & 1/2 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1/3 & 0 \end{bmatrix}$$



$$P = \begin{bmatrix} 1 & 2 & 3 & 4 & 5 \\ 0 & 1/3 & 1/4 & 0 & 0 \\ 1/2 & 0 & 1/4 & 1/3 & 0 \\ 1/2 & 1/3 & 0 & 1/3 & 1/2 \\ 0 & 1/3 & 1/4 & 0 & 1/2 \\ 0 & 0 & 1/4 & 1/3 & 0 \end{bmatrix}$$



- ▶ **Initial state:** The mouse is located at some room  $i$ : probabilities

$$v_0 = (x_1, \dots, x_5).$$

- ▶ Probability mouse starts at room 1 is  $x_1$  item **Transition matrix:**  
 $v_{n+1} = Av_n$  means that  $A$  dictates how probabilities change.
- ▶ Probability mouse is at room 3 after  $n$  steps of the walk:  
 third entry of  $v_n$ .

# Non-regular transition matrix

Disconnected states

Consider the following 'transition graph':



The transition matrix is 
$$\begin{pmatrix} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} \\ 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2} \\ 0 & 0 & \frac{1}{2} & \frac{1}{2} & 0 \end{pmatrix}.$$

Both  $\begin{pmatrix} 1 \\ 1 \\ 0 \\ 0 \\ 0 \end{pmatrix}$  and  $\begin{pmatrix} 0 \\ 0 \\ 1 \\ 1 \\ 1 \end{pmatrix}$ , are eigenvectors with eigenvalue 1.

So there is more than one steady-state vector!

# Stochastic Matrices and Difference Equations

Through an example

Red Box has kiosks all over where you can rent movies. You can return them to any other kiosk.

- ▶ *ij* entry of  $A$ : probability that a movie rented from location  $j$  is returned to location  $i$ .

For example, if there are three locations, maybe

$$A = \begin{pmatrix} .3 & .4 & .5 \\ .3 & .4 & .3 \\ .4 & .2 & .2 \end{pmatrix}.$$

30% probability a movie rented from location 3 gets returned to location 2

**On day  $n$ :**  $x_n, y_n, z_n$  are the numbers of movies in locations 1, 2, 3, respectively, and  $v_n = (x_n, y_n, z_n)$ .

## Probabilistic Intuition

If at opening day the movies are distributed according to  $v_0$  then, **on average:**

$$v_n = Av_{n-1} = A^2v_{n-2} = \cdots = A^n v_0.$$

## Eigenvalues of Stochastic Matrices

**Fact:** 1 is an eigenvalue of a stochastic matrix.

**Why?** If  $A$  is stochastic, then 1 is an eigenvalue of  $A^T$ :

$$\begin{pmatrix} .3 & .3 & .4 \\ .4 & .4 & .2 \\ .5 & .3 & .2 \end{pmatrix} \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix} = 1 \cdot \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}.$$

### Lemma

$A$  and  $A^T$  have the same eigenvalues.

**Proof:**  $\det(A - \lambda I) = \det((A - \lambda I)^T) = \det(A^T - \lambda I)$ , so they have the same characteristic polynomial.

**Note:** This doesn't give a new procedure for finding an eigenvector with eigenvalue 1; it only shows one exists.

**Stronger fact:** if  $\lambda \neq 1$  is an eigenvalue of a *regular* stochastic matrix, then  $|\lambda| < 1$ .

# Eigenvalues of Stochastic Matrices

Continued

So: If  $\lambda$  is an eigenvalue of  $A$  then it is an eigenvalue of  $A^T$ .

$$\text{eigenvector } v = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} \quad \lambda v = A^T v \implies \lambda x_j = \sum_{i=1}^n a_{ij} x_i.$$

*j*th entry of  $A^T v$

Choose  $x_j$  with the largest absolute value, so  $|x_i| \leq |x_j|$  for all  $i$ .

$$|\lambda| \cdot |x_j| = \left| \sum_{i=1}^n a_{ij} x_i \right| \leq \sum_{i=1}^n a_{ij} \cdot |x_i| \leq \sum_{i=1}^n a_{ij} \cdot |x_j| = 1 \cdot |x_j|,$$

*positive*       $= \sum_i a_{ij}$        $\geq |x_j|$

so  $|\lambda| \leq 1$ .

**Conclusion:** if  $\lambda \neq 1$  is an eigenvalue of a stochastic matrix with all entries positive, then  $|\lambda| < 1$ . This proof is adapted for *regular* matrices.

## Dynamical systems

Except for the 1-eigenspace, all the others are shriking!  
What happens in the long run?



# Diagonalizable Stochastic Matrices

Example from §5.3

Let  $A = \begin{pmatrix} 3/4 & 1/4 \\ 1/4 & 3/4 \end{pmatrix}$ . This is a regular stochastic matrix.

We saw last time that  $A$  is diagonalizable:

$$A = PDP^{-1} \quad \text{for} \quad P = \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \quad D = \begin{pmatrix} 1 & 0 \\ 0 & 1/2 \end{pmatrix}.$$

**Change of basis:** Let  $w_1 = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$  and  $w_2 = \begin{pmatrix} 1 \\ -1 \end{pmatrix}$  be the columns of  $P$ .

$$A^n = c_1 w_1 + \frac{c_2}{2^n} w_2.$$

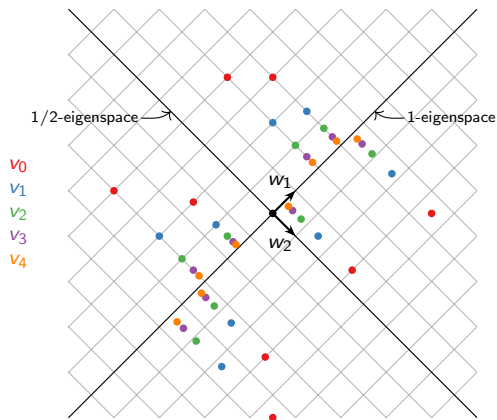
When  $n$  is large, the second term disappears, so  $A^n x$  approaches  $c_1 w_1$ , which is an *eigenvector with eigenvalue 1* (assuming  $c_1 \neq 0$ ).

So all vectors get “sucked into the 1-eigenspace,” which is spanned by  $w_1 = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$ .

# Diagonalizable Stochastic Matrices

Example, continued

Recall:  $A^n = PD^nP^{-1}$  acts on the usual coordinates of  $v_0$  in the same way that  $D^n$  acts on the  $\mathcal{B}$ -coordinates, where  $\mathcal{B} = \{w_1, w_2\}$ .



All vectors get “sucked into the 1-eigenspace.”

## Diagonalizable Stochastic Matrices

The Red Box matrix  $A = \begin{pmatrix} .3 & .4 & .5 \\ .3 & .4 & .3 \\ .4 & .2 & .2 \end{pmatrix}$  is diagonalizable

$$A = P \begin{pmatrix} 1 & 0 & 0 \\ 0 & .1 & 0 \\ 0 & 0 & -.2 \end{pmatrix} P^{-1} = PDP^{-1}.$$

Hence it is easy to compute the powers of  $A$ :

$$A^n = P \begin{pmatrix} 1 & 0 & 0 \\ 0 & (.1)^n & 0 \\ 0 & 0 & (-.2)^n \end{pmatrix} P^{-1} = PD^nP^{-1}.$$

Let  $w_1, w_2, w_3$  be the columns of  $P$ , i.e. the eigenvectors of  $P$  with respective eigenvalues  $1, .1, -.2$ . Let  $\mathcal{B} = \{w_1, w_2, w_3\}$ .

If  $w_1, w_2, w_3$  are the column vectors of  $P$  then

$$A^n x = c_1 w_1 + (.1)^n c_2 w_2 + (-.2)^n c_3 w_3 \rightarrow c_1 w_1$$

# Diagonalizable Stochastic Matrices

## Interpretation

If  $A$  is the Red Box matrix, and  $v_n$  is the vector representing the number of movies in the three locations on day  $n$ , then

$$v_{n+1} = Av_n.$$

For any starting distribution  $v_0$  of videos in red boxes, after enough days, the distribution  $v$  ( $= v_n$  for  $n$  large) is an eigenvector with eigenvalue 1:

$$Av = v.$$

In other words, eventually the number of movies in each kiosk doesn't change much.

Moreover, we know exactly what  $v$  is: a multiple of  $w_1$

- ▶ The entries in  $v$  have to sum up to the number of initial movies (same sum as entries in  $v_0$ ).

(Remember the total number of videos never changes.) Presumably, Red Box

really does have to do this kind of analysis to determine how many videos to put in each box.

## Find the actual Steady State $w_1$

Red Box example

If one computes  $\text{Nul}(A - I)$  and find that  $w' = \begin{pmatrix} 7 \\ 6 \\ 5 \end{pmatrix}$

is an eigenvector with eigenvalue 1.

Then, to get a steady state, divide by  $18 = 7 + 6 + 5$  to get

$$w = \frac{1}{18} \begin{pmatrix} 7 \\ 6 \\ 5 \end{pmatrix} \sim (0.39, 0.33, 0.28).$$

### The long-run

So if you start with 100 total movies, eventually you'll have  $100w = (39, 33, 28)$  movies in the three locations, every day.

### The time spent on a state

Regardless of the initial location of a particular movie. Eventually, that movie will get 'returned' 39% of the times at location 1, 33% at location 2, and 28% at location 3.

## Perron–Frobenius Theorem

These conclusions apply to *any* regular stochastic matrix—whether or not it is diagonalizable!

### Perron–Frobenius Theorem

If  $A$  is a regular stochastic matrix, then it admits a unique steady state vector  $w$ , which spans the 1-eigenspace.

Moreover, for any vector  $v_0$  with entries summing to some number  $c$ , the iterates  $v_1 = Av_0$ ,  $v_2 = Av_1$ ,  $\dots$ ,  $v_n = Av_{n-1}$ ,  $\dots$ , approach  $cw$  as  $n$  gets large.

### Translation:

- ▶ The *1-eigenspace* of a regular stochastic matrix  $A$  *is a line*.
- ▶ The *vector*  $w$  has entries that sum to 1, and are *strictly positive*!
- ▶ Eventually, the *movies arrange* themselves according to the *steady state percentage*, i.e.,  $v_n \rightarrow cw$ .

(The sum  $c$  of the entries of  $v_0$  is the total number of movies)

# Google's PageRank

Internet searching in the 90's was a pain. Yahoo or AltaVista would scan pages for your search text, and just list the results with the most occurrences of those words.

Not surprisingly, the more unsavory websites soon learned that by putting popular words a million times in their pages, they could show up first on popular searches.

Larry Page and Sergey Brin invented a way to rank pages by *importance*. They founded Google based on their algorithm.

Here's how it works. (roughly)

Reference:

<http://www.math.cornell.edu/~mec/Winter2009/RalucaRemus/Lecture3/lecture3.html>

## Google's PageRank: The Importance Rule

Each webpage has an associated importance, or **rank**. This is a positive number.

### The Importance Rule

If page  $P$  *links* to  $n$  other pages  $Q_1, Q_2, \dots, Q_n$ , then each  $Q_i$  *should inherit*  $\frac{1}{n}$  of  $P$ 's *importance*.

- ▶ A very important page links to your webpage: then your webpage is important.
- ▶ A ton of unimportant pages link to your webpage: then it's still important.
- ▶ But if only one crappy site links to yours, your page isn't important.

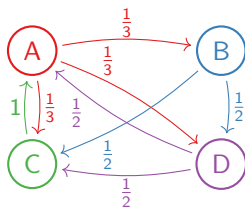
### Random surfer interpretation

A "random surfer" just randomly clicks on link after link. The pages she *spends the most time* on should be *the most important*. **Stochastic terms:** random walk on the graph of hiperlinks. Look for steady-state vector!



# The Importance Matrix

Consider the following Internet with only four pages. Links are indicated by arrows.



In terms of matrices, if  $v = (a, b, c, d)$  is the vector containing the ranks  $a, b, c, d$  of the pages  $A, B, C, D$ , then

**importance matrix:**  $ij$  entry is importance page  $j$  passes to page  $i$

$$\begin{pmatrix} 0 & 0 & 1 & \frac{1}{2} \\ \frac{1}{3} & 0 & 0 & 0 \\ \frac{1}{3} & \frac{1}{2} & 0 & \frac{1}{2} \\ \frac{1}{3} & \frac{1}{2} & 0 & 0 \end{pmatrix} \begin{pmatrix} a \\ b \\ c \\ d \end{pmatrix} = \begin{pmatrix} c + \frac{1}{2}d \\ \frac{1}{3}a \\ \frac{1}{3}a + \frac{1}{2}b + \frac{1}{2}d \\ \frac{1}{3}a + \frac{1}{2}b \end{pmatrix} \stackrel{\text{Importance Rule}}{=} \begin{pmatrix} a \\ b \\ c \\ d \end{pmatrix}$$

# The 25 Billion Dollar Eigenvector

## Observations:

- ▶ The importance matrix is a stochastic matrix!
- ▶ The rank vector is an eigenvector with eigenvalue 1

**Random surfer interpretation:** If a random surfer has probability  $(a, b, c, d)$  to be on page  $A, B, C, D$ , respectively, then after clicking on a random link, the probability he'll be on each page is

$$\begin{pmatrix} 0 & 0 & 1 & \frac{1}{2} \\ \frac{1}{3} & 0 & 0 & 0 \\ \frac{1}{3} & \frac{1}{2} & 0 & \frac{1}{2} \\ \frac{1}{3} & \frac{1}{2} & 0 & 0 \end{pmatrix} \begin{pmatrix} a \\ b \\ c \\ d \end{pmatrix} = \begin{pmatrix} c + \frac{1}{2}d \\ \frac{1}{3}a \\ \frac{1}{3}a + \frac{1}{2}b + \frac{1}{2}d \\ \frac{1}{3}a + \frac{1}{2}b \end{pmatrix}.$$

The rank vector is a *steady state* for the importance matrix : it's the probability vector  $(a, b, c, d)$  such that, after clicking on a random link, the random surfer will have the *same probability* of being on each page.

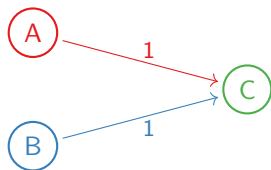
So, the important (high-ranked) pages are those where a random surfer will end up most often.

# Problems with the Importance Matrix

Dangling pages

**Observation:** the importance matrix is *not* regular: it's only nonnegative. So we can't apply the Perron–Frobenius theorem. *How does this cause problems?*

Consider the following Internet:



The importance matrix is  $\begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 1 & 1 & 0 \end{pmatrix}$ . is not stochastic!

$$f(\lambda) = \det \begin{pmatrix} -\lambda & 0 & 0 \\ 0 & -\lambda & 0 \\ 1 & 1 & -\lambda \end{pmatrix} = -\lambda^3.$$

and 1 is not even an eigenvalue: there is no rank vector!

## The Google Matrix (Page and Brin's solution)

Fix  $p$  in  $(0, 1)$ , called the **damping factor**. (A typical value is  $p = 0.15$ .)

The **Google Matrix** is

$$M = (1 - p) \cdot A + p \cdot B \quad \text{where} \quad B = \frac{1}{N} \begin{pmatrix} 1 & 1 & \cdots & 1 \\ 1 & 1 & \cdots & 1 \\ \vdots & \vdots & \ddots & \vdots \\ 1 & 1 & \cdots & 1 \end{pmatrix},$$

$N$  is the total number of pages, and  $A$  is the importance matrix.

- ▶ Random surfer interpretation: with probability  $p$  the surfer gets bored and starts over on a completely random page.

Fact

The PageRank vector is the steady state for the Google Matrix.

This exists and has positive entries by the Perron–Frobenius theorem. The hard part is calculating it: the Google matrix has 1 gazillion rows.